Nonregular Languages - Lecture 7 James Marshall

Examples of nonregular languages (?):

 $C = \{w \mid w \text{ has an equal number of 0s and 1s}\}$

 $D = \{ww^{\mathcal{R}}\}$

An examples that looks nonregular but isn't:

 $E = \{w \mid w \text{ has an equal number of substrings } 01 \text{ and } 10\}$

Exercise: prove the regularity of E

Theorem - The Pumping Lemma for Regular Languages

If *A* is a regular language, then there is a minimum length *p* (the *pumping length*) where, for any string $s \in A$ of length at least *p*, *s* can be divided into three pieces s = xyz such that

1. $xy^i z \in A$, for each $i \ge 0$ 2. |y| > 03. $|xy| \le p$

Proof Sketch

1. Consider a machine M that recognises A

2. Choose a string $s \in A$ of length $n \ge p$, where p = the number of states in M

3. Let $q_1q_2...q_{n+1}$ be the sequence of states *M* goes through in accepting *s*

4. By the *pigeonhole principle*, since $n \ge p$, one or more states in *M* must be visited more than once

5. Consider the first p + 1 states visited by M; one state (labelled q_r) must be visited twice, so *s* can be broken down as follows:

Example 1:

Prove that $F = \{0^{n}1^{n} | n \ge 0\}$ is not regular, by application of the pumping lemma

Example 2:

Prove that $G = \{w \mid w \text{ has an equal number of 0s and 1s}\}$ is not regular, *without* recourse to the pumping lemma